Lezione 14 Bioinformatica

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Lecture 14: Longest Common Subsequence

Longest common subsequence (LCS) problem BLAST (Basic Local Alignment Search Tool) FASTA (FAST Alignement)



Sommario

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Dynamic Programming Approach

LCS : Longest Common Subsequence

Let $X, Y \in Seq$ be the sequences to compare, and X_i, Y_j be the subsequences of their first *i*, *j* characters, respectively. The integer function

$\textit{LCS}:\textit{Seq} \times \textit{Seq} \rightarrow \textit{Nat}$

gives the integer length of longest common subsequence of any two (sub)sequences, as follows:



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gives the integer length of longest common subsequence of any two (sub)sequences, as follows:

$$LCS(X_{i}, Y_{j}) = \begin{cases} 0 & \text{if } i = 0 \text{ or } j = 0 \\ LCS(X_{i-1}, Y_{j-1}) + 1 & \text{if } x_{i} = y_{j} \\ \max(LCS(X_{i}, Y_{j-1}), LCS(X_{i-1}, Y_{j})) & \text{if } x_{i} \neq y_{j} \end{cases}$$



Longest common subsequence

LCS function defined

 $x_i = y_j$

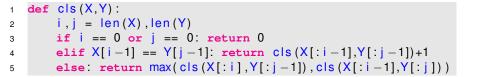
 $LCS(X_i, Y_i) = LCS(X_{i-1}, Y_{i-1}) + 1$ X_i Y_{j-1} X_{i-1} x_i Y_{j-1} y_j X_{i-1} Y_i

 $x_i \neq y_j$ $LCS(X_i, Y_j) = \max(LCS(X_i, Y_{j-1}), LCS(X_{i-1}, Y_j))$



Recursive implementation

just write down in Python the recursive equations above



- 1 print cls("BASKETBALL","BASEBALL") = 8
 OK !
- 1 print cls("ABRACADABRA", "SUPERCALIFRAGILISTICESPIRALIDOSO")
 VERY long execution time ... WHY ?



... because of recursion nonlinearity

the execution time is exponential with the sequence lengths

a recursion is said linear if the definition right-hand side contains at most one recursive function call



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▶ linear recursion:
$$\binom{n}{k} = \binom{n-1}{k-1} \times \frac{n}{k}$$
 complexity: $O(n)$
1 def binomial (n,k) :
2 if k == 0 or n == k: return 1
2 olso: return binomial $(n-1,k-1) \times n / k$



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Memoization technique

In computing, "memoization" is an optimization technique used primarily to speed up computer programs by having function calls avoid repeating the calculation of results for previously-processed input

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- This technique of saving values that have already been calculated is frequently used
- Memoization is a means of lowering a function's time cost in exchange for space cost; that is, memoized functions become optimized for speed in exchange for a higher use of computer memory space.
- An efficient LCS procedure requires: saving the solutions to one level of subproblem in a table so that the solutions are available to the next level of subproblems.



Length of the Longest Common Subsequence

computing the function $LCS: Seq \times Seq \rightarrow Nat$ with memoization

```
def LCS(X, Y):
1
       m,n = len(X), len(Y)
2
       # An (m+1) times (n+1) matrix
3
       C = [[0] * (n+1) \text{ for } i \text{ in } range(m+1)]
4
       for i in range(1, m+1):
5
            for j in range(1, n+1):
6
                if X[i-1] == Y[i-1]:
7
                   C[i][i] = C[i-1][i-1] + 1
8
9
                else
                   C[i][i] = max(C[i][i-1], C[i-1][i])
10
        return C
11
```



```
1 >>> X = "AATCC"

2 >>> Y = "ACACG"

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1	>>> print C								
2	[[0,	0,	0,	0,	0,	0],			
3	[0,	1,	1,	1,	1,	1],			
4	[0,	1,	1,	2,	2,	2],			
5	[0,	1,	1,	2,	2,	2],			
6	[0,	1,	2,	2,	З,	3],			
7	[0,	1,	2,	2,	З,	3]]			



1 >>> X = "ATGGCCTGGAC" 2 >>> Y = "ATCCGGACC" 3 >>> m = len(X) 4 >>> n = len(Y) 5 >>> C = LCS(X, Y)



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2	>>> Y =	"ATCCGGACC"
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4	>>> n =	len(Y)
5	>>> <mark>C</mark> =	LCS(X, Y)

1	>>> [>>> print C								
2	[[0,	0,	0,	0,	0,	0,	0,	0,	0,	0],
3	[0,	1,	1,	1,	1,	1,	1,	1,	1,	1],
4	[0,	1,	2,	2,	2,	2,	2,	2,	2,	2],
5	[0,	1,	2,	2,	2,	З,	З,	З,	З,	3],
6	[0,	1,	2,	2,	2,	З,	4,	4,	4,	4],
7	[0,	1,	2,	З,	З,	З,	4,	4,	5,	5],
8	[0,	1,	2,	З,	4,	4,	4,	4,	5,	6],
9	[0,	1,	2,	З,	4,	4,	4,	4,	5,	6],
10	[0,	1,	2,	З,	4,	5,	5,	5,	5,	6],
11	[0,	1,	2,	З,	4,	5,	6,	6,	6,	6],
12	[0,	1,	2,	З,	4,	5,	6,	7,	7,	7],
13	[0,	1,	2,	З,	4,	5,	6,	7,	8,	8]]



Reading out an LCS

Backtracking on the table from the lower-right corner

```
def backTrack(C, X, Y, i, j):
1
        if i == 0 or i == 0:
2
            return
3
        elif X[i-1] == Y[i-1]:
4
            return backTrack(C, X, Y, i-1, j-1) + X[i-1]
5
        else:
6
            if C[i][i-1] > C[i-1][i]:
7
                 return backTrack(C, X, Y, i, j-1)
8
            else:
9
                 return backTrack(C, X, Y, i-1, j)
10
```



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1 >>> print "Some_LCS:_'%s'" % backTrack(C, X, Y, m, n)
2 Some LCS: 'AAC'



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1 >>> print "Some_LCS:_'%s'" % backTrack(C, X, Y, m, n)
2 Some LCS: 'ATCCGGAC'



Reading out all LCSs

```
def backTrackAll(C, X, Y, i, j):
1
         if i == 0 or i == 0:
2
             return Set([""])
3
         elif X[i-1] == Y[i-1]:
4
             return set ([Z + X[i-1]]
5
                     for Z in backTrackAll(C, X, Y, i-1, j-1)
6
                          ])
         else:
7
             \mathbf{R} = \mathbf{set}()
8
             if C[i][j-1] >= C[i-1][j]:
9
                  R.update(backTrackAll(C, X, Y, i, j-1))
10
             if C[i-1][i] >= C[i][i-1]:
11
                  R.update(backTrackAll(C, X, Y, i-1, j))
12
             return R
13
```



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Sommario

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Comparison of nucleotide or protein sequences

The Basic Local Alignment Search Tool (BLAST) finds regions of local similarity between sequences



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How to do Batch BLAST jobs

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- Using the BLAST web pages it is possible to input "batches" of sequences into one form and retrieve the results
- There are two methods to do batch BLAST jobs
- The first is through the web interface and the second is using the standalone BLAST binaries and downloaded NCBI databases







BLAST paper

QuickStart: Example-Driven Web-Based BLAST Tutorial



Sommario

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Example

FASTA stands for FAST-ALL, reflecting the fact that it can be used for a fast protein comparison or a fast nucleotide comparison

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- The trade-off between speed and sensitivity is controlled by the ktup parameter, which specifies the size of the word
- Increasing the ktup decreases the number of background hits
- Not every word hit is investigated but instead initially looks for segment's containing several nearby hits



FASTA Web services

Both REST and SOAP web service interfaces are exposed

REST Sample clients are provided for a number of programming languages.

SOAP RPC/encoded SOAP service

